

Eigenvalue multiplicity & Equiangular lines

Joint with Alexander Polyanskii
arxiv: 1708.02317.

Jonathan Tidor

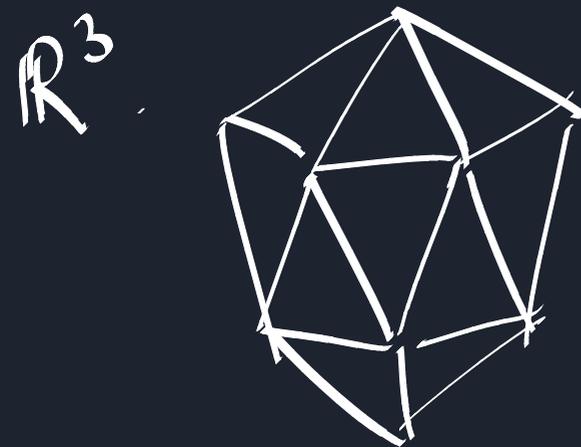
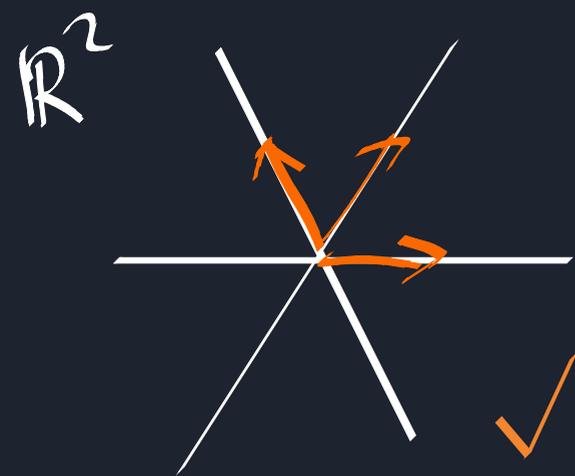
Yuan Yao

Shengtong Zhang

Yufei Zhao

arxiv: 1907.12466

Lines in \mathbb{R}^n (through 0)
pairwise separated by same angle



6 lines.

Question max size of equi.
lines in \mathbb{R}^n ?

Question max size of equi.
lines in \mathbb{R}^n ?

n	2	3-4	5	6	7-14	...
max	3	6	10	16	28	...
					<u>23-41</u>	

$Cn^2 \leq \text{max} \leq \binom{n+1}{2}^{276}.$
 \uparrow de Caen 2000, \uparrow Gerzon 1973

2018, Balla, Dräxler, Sudakov, Keevash
 $E_\alpha(n) \leq 1.93n$ if $n \geq n_0(\alpha)$,
 unless $\alpha = 1/3$

Question What if the angle
is fixed?

$E_\alpha(n) =$ max size of equi.
lines in \mathbb{R}^n with angle
 $= \arccos \alpha.$

1973 Lemmens-Seidel

$$E_{1/3}(n) = 2(n-1) \quad n \geq 15.$$

1989, Neumaier for $n \geq n_0$

$$E_{1/5}(n) = \left\lfloor \frac{3}{2}(n-1) \right\rfloor \checkmark$$

1973, Neumann, $E_\alpha(n) \leq 2n$ unless
 $\alpha = 1/3, 1/5, 1/7, \dots$

2016, Bukh, $E_\alpha(n) \leq C_\alpha \cdot n.$

Conj 1 (Bukh) $E_{1/3}(n) \approx \frac{4}{3}n$

$$E_{\frac{1}{2k-1}}(n) \approx \frac{k}{k-1}n.$$

Conj 2 (J.-Polyanskii).

$$E_{\alpha}(n) \approx \frac{k}{k-1}n, \text{ where}$$

$$k = \underline{k(\lambda)}, \quad \lambda = \frac{1-\alpha}{2\alpha}.$$

Spectral radius order

$k(\lambda) =$ smallest k s.t. \exists

k -vertex graph G s.t. $\lambda_1(G) = \lambda$

$\lambda_1(G) \geq \lambda_2(G) \geq \dots \geq \lambda_k(G)$
eigenvalues of adjacency matrix of G .

α	λ	G	k	$E_{\alpha}(n) \approx$
$1/3$	1	\rightarrow	2	$2n$
$1/5$	2	Δ	3	$\frac{3}{2}n$
$1/7$	3	\boxtimes	4	$\frac{4}{3}n$
$1/(1+2\sqrt{2})$	$\sqrt{2}$	Λ	3	$\frac{3}{2}n$

THM. (JP). Conj 2 holds

for $\lambda \leq \sqrt{2+\sqrt{5}}$.

$\rightarrow \Lambda \Delta \boxtimes$
 $\times \times \times (\times)$
 \uparrow Barrier $1 \quad \sqrt{2} \quad 2 \quad 3R$

THM (JTYZZ).

$$E_\alpha(n) = \left\lfloor \frac{k}{k-1} (n-1) \right\rfloor$$

for all $n \geq n_0(\alpha)$.

where $k = k(\lambda)$, $\lambda = \frac{1-\alpha}{2\alpha}$

Equiangular lines in \mathbb{R}^n .

$V =$ set of unit vectors
(each vector represents a line).

$$\langle v_i, v_j \rangle = \pm \alpha.$$

Gram matrix $(\langle v_i, v_j \rangle)_{i,j} \succeq 0$.

rank (Gram matrix) $\leq n$.

Goal. Given n , find largest m

s.t. \exists m -vertex graph G with

(PSD) + (RANK).

Alternative goal Given m , find

smallest n , s.t. ... (PSD) + (RANK).

\Leftrightarrow minimize $\text{rank}(\lambda I - A)$

m -vertex graph G .

$V =$ vertex set.

$$v_i \sim v_j \Leftrightarrow \langle v_i, v_j \rangle < 0.$$

$$\lambda I - A + \frac{1}{2} J \succeq 0 \quad (\text{PSD})$$

$\lambda = \frac{1-\alpha}{2\alpha}$ \uparrow adj. mat. of G . \leftarrow all-ones mat.

$$\text{rank}(\lambda I - A + \frac{1}{2} J) \leq n \quad (\text{RANK})$$

\Leftrightarrow maximize $\text{mult}(\lambda, A)$.

Alt goal Given m . find

m -vertex G , satisfying

$$\text{(PSD)}: \underbrace{\lambda I - A} + \frac{1}{2} \underbrace{J} \geq 0.$$

that maximizes $\text{mult}(\lambda, A)$

(PSD) $\xrightarrow{\text{Weil's inequality}}$

(Completely reducible):

$G = G_1 \cup \dots \cup G_c$ where
each connected component G_i
satisfies $\lambda_1(G_i) \leq \lambda$.

$$\text{mult}(\lambda, G) = \sum_{i=1}^c \text{mult}(\lambda, G_i) \leq c.$$

Best to choose $|G_i| = k(\lambda)$.
= smallest k s.t. $\exists k$ -vtx G
s.t. $\lambda_1(G) = \lambda$. $\left. \vphantom{\begin{array}{l} \text{Best to choose } |G_i| = k(\lambda). \\ = \text{smallest } k \text{ s.t. } \exists k\text{-vtx } G \\ \text{s.t. } \lambda_1(G) = \lambda. \end{array}} \right\} \frac{m}{k(\lambda)}$

(Irreducible): G is connected,
but $\lambda_2(G) = \lambda$.

$$\text{mult}(\lambda, G) = o(m).$$

Prop. [BDSK], There is

a switching of G s.t.

max deg of G is bounded by
a constant $\Delta = \Delta(\alpha)$.

THM (JTYZZ). For every n -vertex connected graph

G with max deg $\leq \Delta$.

If $\lambda = \lambda_2(G)$, then \dots

$$\text{mult}(\lambda, G) \leq \frac{cn}{\log \log n} = o(n).$$

Perron-Frobenius

Near-miss examples:



Strongly regular graph.



$$\text{mult}(0, G) = \frac{4}{3}$$

LEM1. Every connected

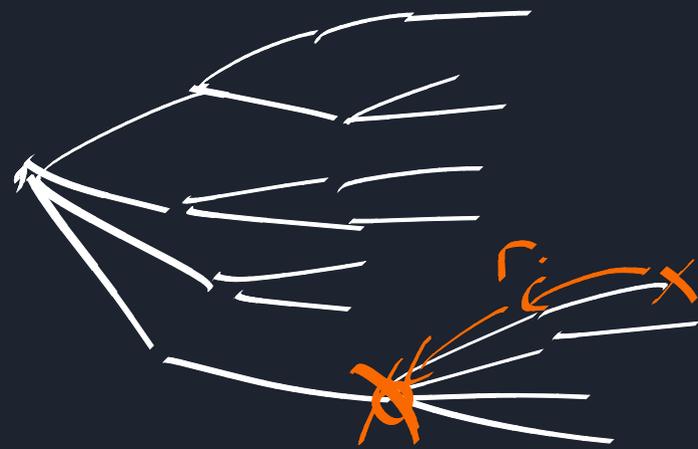
n -vertex graph G has an

r -net of size $\lceil \frac{n}{r+1} \rceil$.

\uparrow \leftarrow vertex set of G .
 $V_0 \subseteq V$ s.t. $\forall v \in V$,

$\exists u \in V_0$. $\text{dist}(u, v) \leq r$.

pf: wlog G is tree.



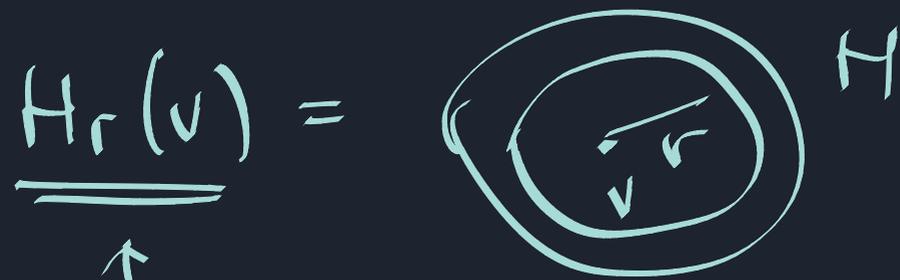
LEM 2. If $H = G - (\text{an } r\text{-net of } G)$.

then $\lambda_1(H)^{2r} \leq \lambda_1(G)^{2r} - 1$.

pf. $A_H^{2r} \leq A_G^{2r} - I_{v \in V_0}$



LEM 3: $\sum_{i=1}^{|H|} \lambda_i(H)^{2r} \leq \sum_{v \in V(H)} \lambda_1(H_r(v))^{2r}$



\uparrow
 r -nbhd of v in H .

pf. LHS = $\text{tr}(A_H^{2r})$.

= $\sum_{v \in V(H)} \#$ of closed walks of length $2r$ starting at v .

$$= \sum_{v \in V(H)} \frac{\mathbf{1}_v^T A_H^{2r} \mathbf{1}_v}{|H_r(v)|}$$

$$\leq \sum_{v \in V(H)} \lambda_1(H_r(v))^{2r} \quad \square$$

pf. $\lambda = \lambda_2(G)$.

$$r = r_1 + r_2, \quad r_1 = c \log \log n$$

$$r_2 = \underline{c} \log n,$$

Case 1. Assume $\exists v$.

$$\lambda_1(G_{r_1}(v)) > \lambda.$$

$$\text{Then } \lambda_1(G - \underline{G_{r_1+1}(v)}) < \lambda.$$

By Cauchy interlacing.

$$\begin{aligned} \text{mult}(\lambda, G) &\leq |G_{r_1+1}(v)| \\ &\leq \frac{\Delta^{r_1+1}}{\lambda} \\ &= o(n). \end{aligned}$$

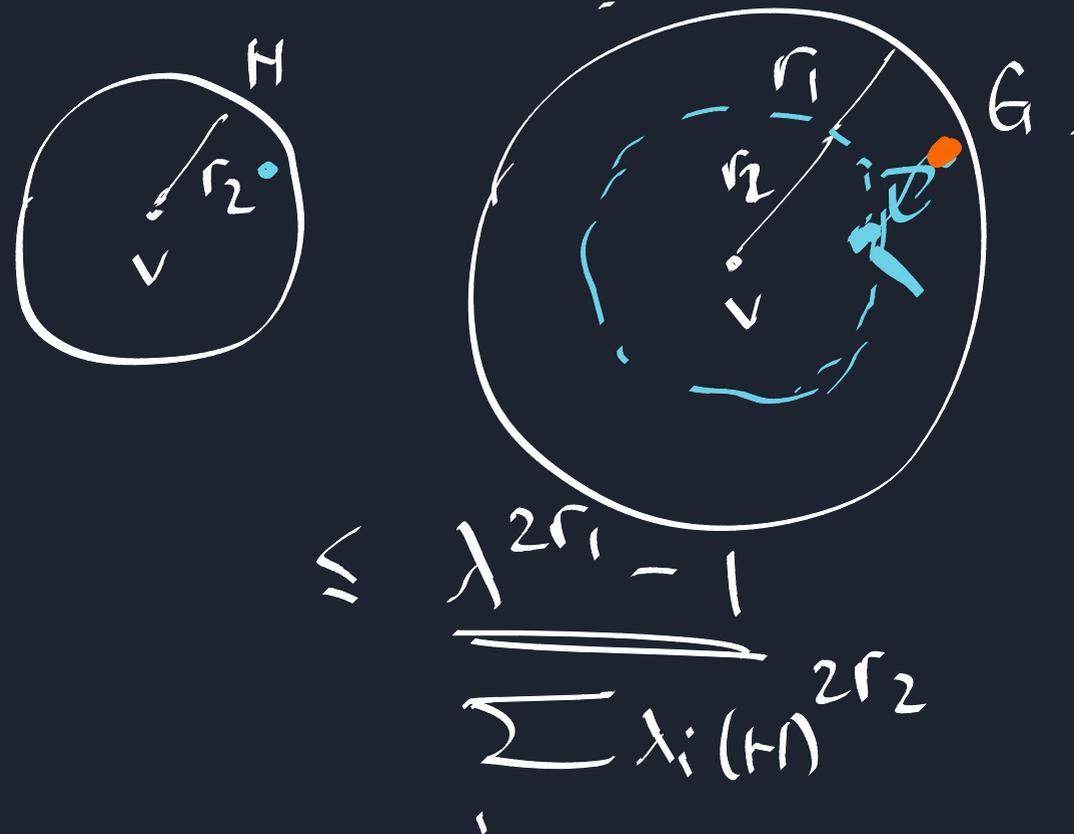
Case 2. Assume $\forall v$.

$$\lambda_1(G_{r_1}(v)) \leq \lambda.$$

Let V_0 be a small r_1 -net of G .

$$H = G - V_0.$$

$$\begin{aligned} \text{LEM 2} &\Rightarrow \lambda_1(H_{r_2}(v))^{2r_1} \\ &\leq \lambda_1(G_{r_1}(v))^{2r_1 - 1}. \end{aligned}$$



$$\text{LEM 3} \Rightarrow \text{mult}(\lambda, H) \cdot \lambda^{2r_2}$$

$$\leq \sum_{i=1}^{|H|} \lambda_i (H)^{2r_2}$$

$$\leq \sum_{v \in V(H)} \lambda_i (H_{r_2}(v))^{2r_2}$$

$$\leq |H| \cdot \left(\lambda^{2r_1} - 1 \right)^{r_2/r_1}$$

$$\Rightarrow \text{mult}(\lambda, H) = o(n).$$

Cauchy interlacing

$$\Rightarrow \text{mult}(\lambda, G) \leq \text{mult}(\lambda, H)$$

$$+ |V_0| = o(n).$$